

# Linear Programming Duality

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## 1 Definition

Let  $A \in \mathbb{R}^{m \times n}$ ,  $I \subseteq [m]$ ,  $J \subseteq [n]$ . Define the linear program  $P := \text{LP}(A, b, c, I, J)$  as

$$\min_{x \in \mathbb{R}^n: x_J \geq 0} c^T x \quad \text{where} \quad ((Ax)_i \geq b_i, \forall i \in I) \text{ and } ((Ax)_i = b_i, \forall i \in [m] - I)$$

Then the dual of  $P$  is  $D := \text{LP}(-A^T, -c, -b, J, I)$ , i.e.,

$$\max_{y \in \mathbb{R}^m: y_I \geq 0} b^T y \quad \text{where} \quad ((A^T y)_j \leq c_j, \forall j \in J) \text{ and } ((A^T y)_j = c_j, \forall j \in [n] - J)$$

**Lemma 1.** *Dual of dual is dual.*

On setting  $J = [n]$  and  $I = [m]$ , we get that the following are duals of each other

$$P : \min_{x \geq 0} c^T x \text{ where } Ax \geq b \qquad D : \max_{y \geq 0} b^T y \text{ where } A^T y \leq c$$

## 2 Method for Quick Application

Adapted from Sébastien Lahaie's [notes](#).

1. Express problem in standard form:

- (a) Express as a minimization problem.
- (b) Write each (non-trivial) constraint as  $f(x) \leq 0$  or  $f(x) = 0$ .
- (c) Write each trivial constraint as  $x \geq 0$  (i.e., if  $x \leq 0$ , replace  $x$  by  $-x$ ):

$$\min_{x \in \mathbb{R}^n: x_J \geq 0} c^T x \quad \text{where} \quad (b_i - (Ax)_i \leq 0, \forall i \in I) \text{ and } (b_i - (Ax)_i = 0, \forall i \in [m] - I).$$

2. Add dual variables:

- (a) Create a non-negative dual variable  $y_i$  for each inequality constraint  $f_i(x) \leq 0$ .
- (b) Create an unrestricted dual variable  $y_i$  for each equality constraint  $f_i(x) = 0$ .
- (c) Remove the constraint and add the term  $y_i f_i(x)$  to the objective.
- (d) Maximize over dual variables.

$$\max_{y \in \mathbb{R}^m: y_I \geq 0} \min_{x \in \mathbb{R}^n: x_J \geq 0} c^T x + y^T (b - Ax).$$

3. Rearrange terms to express objective as an affine function of primal variables:

$$\max_{y \in \mathbb{R}^m: y_I \geq 0} \min_{x \in \mathbb{R}^n: x_J \geq 0} b^T y + x^T (c - A^T y).$$

4. For each term  $x_j g_j(y)$  in the objective, remove the term and add constraint

(a)  $g_j(y) \geq 0$  if  $x_j$  is non-negative.

(b)  $g_j(y) = 0$  if  $x_j$  is unrestricted.

$$\max_{y \in \mathbb{R}^m: y_J \geq 0} b^T y \quad \text{where} \quad (c_j - (A^T y)_j \geq 0, \forall j \in J) \quad \text{and} \quad (c_j - (A^T y)_j = 0, \forall j \in [n] - J).$$

5. Rearrange into suitable form

$$\max_{y \in \mathbb{R}^m: y_J \geq 0} b^T y \quad \text{where} \quad ((A^T y)_j \leq c_j, \forall j \in J) \quad \text{and} \quad ((A^T y)_j = c_j, \forall j \in [n] - J).$$