Stochastic Processes

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Definition 1 (Stochastic Process). Let $\mathcal{T} \subseteq \mathbb{R}$. For any $t \in \mathcal{T}$, let X_t (or X(t)) be a random variable with support D. Then $X := \{X_t : t \in \mathcal{T}\}$ is called a stochastic process on state-space D and time \mathcal{T} . Usually, \mathcal{T} is either $\mathbb{Z}_{\geq 0}$ (discrete-time) or $\mathbb{R}_{\geq 0}$ (continuous-time).

1 Discrete-Time Markov Chains

Definition 2 (Markov Chain). Let $X := [X_0, X_1, \ldots]$ be a stochastic process on statespace D and time $\mathbb{Z}_{\geq 0}$. X is called a discrete-time markov chain if $\Pr(X_{t+1} = d \mid X_t, X_{t-1}, \ldots, X_0) = \Pr(X_{t+1} = d \mid X_t)$. If $\Pr(X_{t+1} = v \mid X_t = u) = \Pr(X_1 = v \mid X_0 = u)$ for all t, u, v, then X is called time-homogeneous.

Definition 3 (Transition function). Let X be a markov chain on state space D. Define $P^{(k)}: D \times D \mapsto [0,1]$ as $P^{(k)}(i,j) = \Pr(X_k = j \mid X_0 = i)$. Then $P^{(k)}$ is called the k-step transition function of X. For k = 1, we simply write P instead of $P^{(1)}$. For a finite state space, we can represent P as a matrix.

Lemma 1 (Chapman-Kolmogorov Equation). $P^{(m+n)}(i,j) = \sum_k P^{(m)}(i,k)P^{(n)}(k,j)$.

1.1 Classification of States, Recurrence, Limiting Probabilities

Definition 4. Let $f_{i,j} := \Pr\left(\bigvee_{t\geq 1}(X_t=j) \mid X_0=i\right)$. Then $f_{i,j}$ is called the eventual transition probability from *i* to *j*. If i = j, then we write $f_{i,i}$ as f_i , and call it the recurrence probability of state *i*.

Definition 5. For a state *i*, let N_i be the random variable that counts the number of times we are in state *i*, i.e., $N_i := \sum_{t=0}^{\infty} \mathbf{1}(X_t = i)$. Then N_i is called the visit-count of *i*.

Definition 6. A state *i* of a markov chain is recurrent iff (the following are equivalent):

- the recurrence probability (f_i) of i is 1.
- *i* is visited infinitely often, i.e., $Pr(N_i = \infty | X_0 = i) = 1$.
- *i* is visited infinitely often in expectation, i.e., $E(N_i | X_0 = i) = \infty$.

A non-recurrent state is called a transient state.

Lemma 2. $\Pr(N_i = k \mid X_0 = i) = f_i^{k-1}(1 - f_i).$

Lemma 3. $E(N_i \mid X_0 = i) = 1/(1 - f_i) = \sum_{t=0}^{\infty} P^{(t)}(i, i).$

Definition 7. State j is accessible from state i if $P^{(t)}(i, j) > 0$ for some t. States i and j communicate (denoted as $i \leftrightarrow j$) if i and j are both accessible from each other.

Lemma 4. Accessibility is reflexive and transitive. Communication is an equivalence relation. The equivalence classes of communicability are called state classes. A markov chain is irreducible if it has just one state class.

Definition 8. Let T_i be the time when a markov chain moves to state i, i.e., $T_i := \min_{t \ge 1} (X_t = i)$. When conditioned on $X_0 = i$, T_i is called the recurrence time of i. State i is called positive recurrent if $E(T_i | X_0 = i)$ is finite, otherwise it is null recurrent.

Lemma 5. Recurrence and positive recurrence are class properties, i.e., they are same for all states in a class.

Lemma 6. In a finite-state markov chain, all recurrent states are positive recurrent, and there is at least one recurrent state.

Definition 9 (Periodicity). For a state *i*, its period is defined as $gcd(\{t : Pr(T_i = t \mid X_0 = i) > 0\})$. A state is aperiodic if its period is 1.

Lemma 7. Periodicity is a class property.

Definition 10 (Ergodicity). A state is ergodic if it is positive recurrent and aperiodic. A markov chain is ergodic if all its states are ergodic.

Lemma 8. In an irreducible ergodic markov chain, for every state j, $\lim_{t\to\infty} P^{(t)}(j,i) = \pi_i$ for a unique real number π_i . π_i is called the limiting probability of state i. Furthermore, π_i is the unique solution to this system of equations: $\pi_i = \sum_j \pi_j P(j,i)$ for all i ($\pi = P^T \pi$ in matrix form) and $\sum_i \pi_i = 1$.

Lemma 9. In an irreducible ergodic markov chain, $E(T_i | X_0 = i) = 1/\pi_i$.

Corollary 9.1. A state *i* is null recurrent iff $\pi_i = 0$.

Theorem 10. If the transition function of markov chain X is doubly-stochastic (i.e., each row and each column sums to 1), then the limiting probability of each state is 1/n, where n is the number of states.

1.2 Time-Reversibility

Definition 11. For an irreducible ergodic markov chain X with limiting probabilities π . Let Y be a markov chain whose transition function is $Q(i, j) = P(j, i)(\pi_j/\pi_i)$. Then Y is called the time-reversed markov chain of X. X is called time-reversible if Q = P.

Theorem 11. Let X be a time-reversible markov chain with limiting probabilities π . Then π is the unique solution to this system of equations: $x_j P(j,i) = x_i P(i,j)$ for all states i and j, and $\sum_i x_i = 1$.

Theorem 12. If the transition function of markov chain X is symmetric, then X is time-reversible.

1.3 Simple Random Walk

Let X be a TH MC on state space $S = I \cap \mathbb{Z}$, where I is an interval of \mathbb{R} and

$$\Pr(X_1 = j \mid X_0 = i) = \begin{cases} p & \text{if } j = i+1\\ 1-p & \text{if } j = i-1\\ 0 & \text{if } j \neq i \end{cases}$$

Lemma 13 (Martingale property). Let $\forall t \geq 1$. If p = 1/2, then $E(X_t \mid X_{t-1}) = X_{t-1}$. Otherwise, for r := (1-p)/p, $E(r^{X_t} \mid X_{t-1}) = r^{X_{t-1}}$.

Lemma 14. Let I = [0, b], $T := \min_{t \ge 0} (X_t \in \{0, b\})$, and r := (1 - p)/p. Then $\forall i \in S$,

$$p_i := \Pr(X_T = b \mid X_0 = i) = \begin{cases} \frac{i}{b} & \text{if } p = 1/2 \\ \frac{r^i - 1}{r^b - 1} & \text{if } p \neq 1/2 \end{cases}$$
$$\mu_i := \operatorname{E}(T \mid X_0 = i) = \begin{cases} i(b - i) & \text{if } p = 1/2 \\ \frac{1}{1 - 2p} \left(i - b\frac{r^i - 1}{r^b - 1}\right) & \text{if } p \neq 1/2 \end{cases}$$

Proof sketch. $p_0 = 0$, $p_b = 1$, and $p_i := pp_{i+1} + (1-p)p_{i-1} \quad \forall i \in [b-1]$. Rearrange to get $p(p_{i+1} - p_i) = (1-p)(p_i - p_{i-1}) \quad \forall i \in [b-1]$. Let $d_i := p_i - p_{i-1} \quad \forall i \in [b]$. Then

$$\begin{aligned} \forall i \in [b], \quad \frac{d_i}{d_1} &= \prod_{j=1}^{i-1} \frac{d_{j+1}}{d_j} = \prod_{j=1}^{i-1} r = r^{i-1}, \\ \forall i \in S, \quad p_i &= \sum_{j=1}^i d_j = \sum_{j=1}^i d_1 r^{j-1} = d_1 \times \begin{cases} i & \text{if } p = 1/2 \\ \frac{r^i - 1}{r - 1} & \text{if } p \neq 1/2 \end{cases}. \end{aligned}$$

Since $p_b = 1$, we get $d_1 = 1/b$ if p = 1/2 and $d_1 = (r-1)/(r^b - 1)$ otherwise.

 $\mu_0 = \mu_b = 0$, and $\mu_i = 1 + p\mu_{i+1} + (1-p)\mu_{i-1} \quad \forall i \in [b-1]$. Rearrange to get $p(\mu_{i+1} - \mu_i) = (1-p)(\mu_i - \mu_{i-1}) - 1 \quad \forall i \in [b-1]$. Let $\nu_i := \mu_i - \mu_{i-1} \quad \forall i \in [b]$. Then $\forall i \in [b-1], \ \nu_{i+1} = r\nu_i - 1/p$.

Case 1: p = 1/2

$$\forall i \in [b], \quad \nu_i - \nu_1 = \sum_{j=1}^{i-1} (\nu_{j+1} - \nu_j) = -2(i-1).$$

$$\forall i \in S, \quad \mu_i = \sum_{j=1}^i \nu_j = \sum_{j=1}^i (\nu_1 - 2(j-1)) = i(\nu_1 - (i-1)).$$

Since $\mu_b = 0$, we get $\nu_1 = b - 1$, and so $\mu_i = i(b - i)$.

Case 2: $p \neq 1/2$

$$\nu_{i+1} = r\nu_i - \frac{1}{p} = r\nu_i - \frac{1}{p}\left(\frac{r}{r-1} - \frac{1}{r-1}\right) \implies \nu_{i+1} - \frac{1}{p(r-1)} = r\left(\nu_i - \frac{1}{p(r-1)}\right).$$

p(r-1) = 1 - 2p. Hence,

$$\begin{aligned} \forall i \in [b], \quad \frac{\nu_i - \frac{1}{1 - 2p}}{\nu_1 - \frac{1}{1 - 2p}} &= \prod_{j=1}^{i-1} \frac{\nu_{j+1} - \frac{1}{1 - 2p}}{\nu_j - \frac{1}{1 - 2p}} = r^{i-1}. \end{aligned}$$
$$\forall i \in S, \quad \mu_i = \sum_{j=1}^i \nu_i = \sum_{j=1}^i \left(\frac{1}{1 - 2p} + r^{j-1} \left(\nu_1 - \frac{1}{1 - 2p}\right)\right). \\ &= \frac{i}{1 - 2p} + \left(\nu_1 - \frac{1}{1 - 2p}\right) \frac{r^i - 1}{r - 1}. \end{aligned}$$

Set $\mu_b = 0$ to get the answer.

Lemma 15 (Catalan number). The number of balanced parentheses strings of length 2n is $\binom{2n}{n}\frac{1}{n+1}$.

Lemma 16. For all $n \geq 1$,

$$\binom{2n}{n}\frac{\sqrt{n}}{4^n} \in \left[\frac{2\sqrt{\pi}}{e^2}, \frac{e}{\sqrt{2\pi}}\right].$$

Proof. Use Stirling's approximation: $n!e^n/(n^n\sqrt{n}) \in [\sqrt{2\pi}, e]$.

Lemma 17. Let $S = \mathbb{Z}$. If $p \neq 1/2$, then every state is transient. Otherwise, every state is null recurrent.

Proof. All states communicate, so all states belong to the same class. Recurrence and positive recurrence are class properties.

Using Definition 6, we get that state 0 is recurrent iff $E(N_0 | X_0 = 0)$, where N_0 is the number of times we visit state 0.

$$E(N_0 \mid X_0 = 0) = \sum_{t=0}^{\infty} P^{(t)}(0, 0)$$
 (by Lemma 2)

$$=\sum_{j=0}^{\infty} 2\binom{2j}{j} p^{j} (1-p)^{j}$$
(by Lemma 15)

$$\leq \frac{\sqrt{2}e}{\pi} \sum_{j=0}^{\infty} \frac{(4p(1-p))^j}{\sqrt{j}}.$$
 (by Lemma 16)

 $4p(1-p) \le 1$ and $p = 1/2 \iff 4p(1-p) = 1$. Hence, the series is convergent iff $p \ne 1/2$. Hence, state 0 is recurrent iff p = 1/2.

Let p = 1/2. Then $T := \min_{t \ge 0} (X_t = 0)$ and $T' := \min_{t \ge 1} (X_t = 0)$. Then state 0 is null recurrent iff $E(T' \mid X_0 = 0) = \infty$. Since state 0 is recurrent, $Pr(T' = \infty \mid X_0 = 0) = \infty$

0. Hence, $0 = \Pr(T' = \infty \mid X_0 = 0) = \Pr(T = \infty \mid X_0 = 1).$

$$E(T' \mid X_0 = 0) = 1 + E(T \mid X_0 = 1)$$

= $1 + \sum_{j=0}^{\infty} (2j+1) \Pr(T = 2j+1 \mid X_0 = 1)$ (since $\Pr(T = \infty \mid X_0 = 1) = 0$)
= $1 + \sum_{j=0}^{\infty} \frac{2j+1}{j+1} {2j \choose j} \frac{1}{2^{2j+1}}$ (by Lemma 15)
 $\ge 1 + \frac{\sqrt{\pi}}{e^2} \sum_{j=0}^{\infty} \frac{1}{\sqrt{j}}$ (by Lemma 16)
= ∞ .

Hence, 0 is null recurrent.

Lemma 18. Let $I = [0, \infty)$. Let $T := \min_{t \ge 0} (X_t = 0)$. Then for i > 0,

$$\mu_i := \mathcal{E}(T \mid X_0 = i) = \begin{cases} i/(1-2p) & \text{if } p < 1/2\\ \infty & \text{if } p \ge 1/2 \end{cases}$$
$$p_i := \Pr(T \neq \infty \mid X_0 = i) = \begin{cases} 1 & \text{if } p \le 1/2\\ r^i & \text{if } p > 1/2 \end{cases}$$

Proof sketch. $\mu_i = i\mu_1$ and $\mu_1 = 1 + p\mu_2 = 1 + 2p\mu_1$. Suppose $\mu_1 \neq \infty$. If p > 1/2, then $\mu_1 = 1/(1-2p) < 0$, which is a contradiction. If p = 1/2, then $\mu_1 = 1 + \mu_1$, which is a contradiction. Hence, $\mu_1 = \infty$ when $p \ge 1/2$.

 $p_i = p_1^i$ and $p_1 = (1 - p) + pp_2$. Solving these equations gives us $p_1 \in \{1, r\}$. When $p \leq 1/2$, we get $r \geq 1$, but $p_1 \in [0, 1]$ (since p_1 is a probability). Hence, $p_i = 1$ when $p \leq 1/2$.

Let p < 1/2. Then $\mu_1 \in \{\infty, 1/(1-2p)\}$. We will show that $\mu_1 \neq \infty$. Since $p_i = 1$ for all $i \ge 0$, we have $\Pr(T = \infty \mid X_0 = i) = 0$. Hence,

$$\mu_{1} = \mathcal{E}(T \mid X_{0} = 1) = \sum_{j=0}^{\infty} (2j+1) \operatorname{Pr}(T = 2j+1 \mid X_{0} = 1)$$

$$= \sum_{j=0}^{\infty} \frac{2j+1}{j+1} {2j \choose j} p^{j} (1-p)^{j+1} \qquad \text{(by Lemma 15)}$$

$$\leq 2(1-p) \frac{e}{\sqrt{2\pi}} \sum_{j=0}^{\infty} \frac{(4p(1-p))^{j}}{\sqrt{j}}. \qquad \text{(by Lemma 16)}$$

This series is convergent, since 4p(1-p) < 1 for p < 1/2. Hence, $\mu_1 \neq \infty$, so $\mu_1 = 1/(1-2p)$.

Let p > 1/2. Consider the random walk Y on state space \mathbb{Z} where $\forall t \ge 0$, $\Pr(Y_{t+1} = j+1 \mid Y_t = j) = p$ and $\Pr(Y_{t+1} = j-1 \mid Y_t = j) = 1-p$. Let $T' := \min_{t\ge 0}(Y_t = 0)$. Then $\Pr(T \ne \infty \mid X_0 = 1) = 1 \iff \Pr(T' \ne \infty \mid Y_0 = 1) = 1$. $\Pr(T' \ne \infty \mid Y_0 = -1) = 1$, using the p < 1/2 case. Hence, we return to state 0 in Y with probability 1. Hence, 0 is a recurrent state in Y, which is a contradiction. Hence, $p_1 = \Pr(T \ne \infty \mid X_0 = 1) \ne 1$. Hence, $p_1 = r$.

2 Counting Process

Definition 12 (Counting Process). Let N be a stochastic process on state space $\mathbb{Z}_{\geq 0}$ and time $\mathbb{R}_{\geq 0}$. Then N is called a counting process if N(0) = 0 and N(t) is monotone in t, i.e., $t_1 < t_2 \implies N(t_1) \leq N(t_2)$.

Definition 13 (Independent increments). A counting process N has independent increments iff for any two disjoint intervals $(u_1, v_1]$ and $(u_2, v_2]$ in $\mathbb{R}_{\geq 0}$, the random variables $N(v_1) - N(u_1)$ and $N(v_2) - N(u_2)$ are independent.

Definition 14 (Stationary increments). A counting process N has stationary increments iff for any $u \leq v$, the random variables N(v) - N(u) and N(v - u) have the same distribution.

Definition 15 (Arrival and interarrival times). For a counting process N, for $i \in \mathbb{Z}_{\geq 0}$, define the i^{th} arrival time $S_i := \min_{t \geq 0} (N(t) = i)$. For $i \in \mathbb{Z}_{\geq 1}$, define the i^{th} interarrival time $T_i := S_i - S_{i-1}$.

Lemma 19. For a counting process N with arrival times S, $N(t) \ge n \iff S_n \le t$.

Definition 16 (Stopping time). Let $X = [X_1, X_2, ...]$ be a sequence of random variables. The random variable N is called a stopping time for X if for all $n \ge 0$, (the following two definitions are equivalent):

- N = n is independent of X_{n+1}, X_{n+2}, \ldots
- $N \leq n$ is independent of X_{n+1}, X_{n+2}, \ldots

Theorem 20 (Wald's identity). Let $X = [X_1, X_2, ...]$ be a sequence of random variables where $E(X_i) = \mu$ for all *i*. Let N be a stopping time for X. Then

$$\operatorname{E}\left(\sum_{i=1}^{N} X_{i}\right) = \mu \operatorname{E}(N).$$

Proof sketch. For all $i, N \ge i$ is independent of X_i , and $\sum_{i=1}^N X_i = \sum_{i=1}^\infty X_i \mathbf{1}(N \ge i)$. \Box

3 Poisson Process

Definition 17 (Poisson process). A counting process N is a Poisson process with rate function $\lambda : \mathbb{R}_{\geq 0} \to \mathbb{R}_{\geq 0}$ if N has independent increments and $N(t_2) - N(t_1) \sim \text{Poisson}(\mu)$, where $\mu := \int_{t_1}^{t_2} \lambda(t) dt$. N is called homogeneous if $\lambda(t) = \lambda(0)$ for all t, otherwise it is called inhomogeneous. For a homogeneous process, we denote $\lambda(0)$ by λ .

Lemma 21. A Poisson process N is homogeneous iff it has stationary increments.

Theorem 22 (Alternative definition of Poisson process). A counting process N is a Poisson process with continuous rate function λ iff N has independent and stationary increments and $\Pr(N(t+h)-N(t)=1) = \lambda(t)h+o(h)$ and $\Pr(N(t+h)-N(t) \ge 2) = o(h)$. Proof sketch for homogeneous. Let $g(u,t) := \mathrm{MGF}_u(N(t)) = \mathrm{E}(e^{uN(t)})$. Show $g(u,t) = 1 + \lambda t(e^u - 1) + o(t)$ straightforwardly. Use calculus to show that $g(u,t) = \exp(e^{\lambda t}(e^u - 1))$ (find derivative w.r.t t by computing $\lim_{h\to 0} (g(u,t+h) - g(u,t))/h$; this gets rid of o(h)). Conclude that $N(t) \sim \mathrm{Poisson}(\lambda t)$ since g(u,t) is MGF of $\mathrm{Poisson}(\lambda t)$. \Box

Lemma 23. For a homogeneous Poisson process N,

$$\Pr(N(s) = a \mid N(s+t) = a+b) = \binom{a+b}{a} \left(\frac{s}{s+t}\right)^a \left(\frac{t}{s+t}\right)^b.$$

Theorem 24. Let N be a counting process. Then N is a homogeneous Poisson process with rate λ iff all interarrival times are independent and distributed $\text{Expo}(\lambda)$.

Theorem 25 (Decomposition theorem 1). Let K be a finite set, and let $\{N_i : i \in K\}$ be independent Poisson processes, where N_i has rate function λ_i . Let $N := \sum_{i \in K} N_i$. Then N is a Poisson process with rate function $\sum_{i \in K} \lambda_i$.

Theorem 26 (Decomposition theorem 2). Let N be a Poisson process with rate function λ . Let K be a finite set (called set of labels). Suppose the j^{th} event receives label $L_j \in K$, where $\Pr(L_j = i) = p_i(S_j)$ for some function $p_i : \mathbb{R}_{\geq 0} \mapsto \mathbb{R}_{\geq 0}$, and $\{N, L_1, L_2, \ldots\}$ are independent. For $i \in K$, let $N_i(t)$ be the number of events having label i, i.e., $N_i(t) = \sum_{j=1}^{N(t)} \mathbf{1}(L_j = i)$. Then N_i is a Poisson process with rate function $p_i\lambda$. Furthermore, all N_i are independent and if all p_i are constant, then $N_i(t) \mid N(t) \sim \operatorname{Binom}(N(t), p_i)$.

Lemma 27. Let $N^{(1)}$ and $N^{(2)}$ be independent homogeneous Poisson processes with rates λ_1 and λ_2 . Then

$$\Pr(S_n^{(1)} < S_m^{(2)}) = \sum_{i=n}^{n+m-1} \binom{n+m-1}{i} \frac{\lambda_1^i \lambda_2^{n+m-1-i}}{(\lambda_1 + \lambda_2)^{n+m-1}}.$$

Proof sketch. Model as a continuous markov chain with state space (n_1, n_2) , where n_i is the number of events of $N^{(i)}$ that have occurred.

Theorem 28 (arrival times distributed as order statistics). Let $X = [X_1, X_2, ..., X_n]$ be IID uniform variables over [0,t]. Let U = sorted(X). Let N be a homogeneous Poisson process. Let S_i be the *i*th arrival time of N. Then conditioned on N(t) = n, the distribution of $[S_1, ..., S_n]$ and U are identical.

Lemma 29 (Excess and Residual). Let N be a Poisson process with rate function λ . Let S_i be the *i*th arrival time. Let $Y(t) := S_{N(t)+1} - t$ and $R(t) := t - S_{N(t)}$. Then $Y(t) > s \iff N(t+s) - N(t) = 0$ and $R(t) > r \iff N(t) - N(t-r) = 0$. If N is homogeneous, we get $Y(t) \sim \text{Expo}(\lambda)$ and $R(t) \sim \text{Expo}(\lambda)$.

4 Continuous-Time Markov Chain

Definition 18 (CTMC). Let $X := \{X(t) : t \in \mathbb{R}_{\geq 0}\}$ be a stochastic process on discrete state-space D. X is called a continuous-time markov chain (CTMC) if $\Pr(X(t+s) = d \mid \{X(u) : 0 \le u \le s\}) = \Pr(X(t+s) = d \mid X(s))$ for all $s, t \in \mathbb{R}_{\geq 0}$. If $\Pr(X(t+s) = v \mid X(s) = u) = \Pr(X(t) = v \mid X(0) = u)$ for all u, v, s, t, then X is called time-homogeneous (TH) or stationary. **Theorem 30** (Equiv defn of TH CTMC). Let $X := \{X(t) : t \in \mathbb{R}_{\geq 0}\}$ be a stochastic process on discrete state-space D. Let $Y(t) := \{X(u) : 0 \leq u < t\}$. Let $T_i^{(s)} := \min_{t \geq 0} (X(t+s) \neq i)$. Let $P_{i,j}^{(s)} := \Pr(X(s+T_i^{(s)}) = j \mid X(s) = i, Y(s))$. X is TH CTMC iff $(T_i^{(s)} \mid X(s) = i, Y(s)) \sim \operatorname{Expo}(\nu_i)$, where ν_i is a constant that doesn't depend on s or Y(s), and $P_{i,j}^{(s)}$ is a constant that doesn't depend on s or Y(s).

Since $T_i^{(s)}$ and $P_{i,j}^{(s)}$ don't depend on s, we simply write T_i and $P_{i,j}$. T_i is called the transition time out of state i, ν_i is called the transition rate out of state i, and $P_{i,j}$ is the probability of transitioning from state i to state j.

Let $q_{i,j} := \nu_i P_{i,j}$. Then $\nu_i = \sum_j q_{i,j}$.

Theorem 31 (Chapman-Kolmogorov DiffEqs). For a TH CTMC X, let $P_{i,j}(t) := \Pr(X(t) = j \mid X(0) = i)$. Then

• Backward DiffEqs:
$$\frac{dP_{i,j}(t)}{dt} = \sum_{k \neq i} q_{i,k} P_{k,j}(t) - \nu_i P_{i,j}(t).$$

• Forward DiffEqs:
$$\frac{dP_{i,j}(t)}{dt} = \sum_{k \neq j} P_{i,k}(t)q_{k,j} - P_{i,j}(t)\nu_j$$
.

Let $r_{i,j} := \begin{cases} q_{i,j} & \text{if } i \neq j \\ -\nu_i & \text{if } i = j \end{cases}$. Let the state space be [n]. Let R be a matrix where $R[i, j] = r_{i,j}$. Then CBKE becomes P'(t) = RP(t) and CFKE becomes P'(t) = P(t)R.

Lemma 32. CKBE P'(t) = RP(t) solves to $P(t) = e^{Rt}$, where $e^A := \sum_{i=0}^{\infty} A^i/i!$ for any square matrix A. Suppose R has n eigenpairs $\{(\lambda_1, v_i) : i \in [n]\}$. Let P be a square matrix whose i^{th} column is v_i , and D be a diagonal matrix whose i^{th} diagonal entry is λ_i . Then $R = PDP^{-1}$, $e^{Rt} = Pe^{Dt}P^{-1}$, and $e^{Dt} = \text{diag}([e^{\lambda_1 t}, \dots, e^{\lambda_n t}])$.

Lemma 33. Let X be a TH CTMC.

$$\lim_{h \to 0} \frac{1 - P_{i,i}(h)}{h} = \nu_i \quad \forall i \qquad \qquad \lim_{h \to 0} \frac{P_{i,j}(h)}{h} = q_{i,j} \quad \forall i \neq j$$

Lemma 34 (Limiting probability). In an irreducible positive-recurrent TH CTMC X, for every state j, $\lim_{t\to\infty} P_{j,i}(t) = P_i$ for a unique real number P_i . P_i is called the limiting probability of state i. Furthermore, P_i is the unique solution to $\sum_i P_i = 1$ and CK forward equations, i.e., $P_i\nu_i = \sum_{j\neq i} P_j q_{j,i}$.

Lemma 35 (Limiting probability of embedded chain). Let X be an irreducible positiverecurrent TH CTMC. Let Y be the sequence of states visited by X. Then Y is a discrete MC. Let P and π be the limiting probabilities of X and Y, respectively. Then $P_i = (\pi_i/\nu_i)/(\sum_j \pi_j/\nu_j)$ and $\pi_i = P_i\nu_i/(\sum_j P_j\nu_j)$.

Definition 19. A CTMC is time-reversible iff the corresponding embedded discrete-time MC is time-reversible.

Lemma 36 (2-state). For a CTMC on states $\{0,1\}$, where $q_{0,1} = \lambda$ and $q_{1,0} = \mu$, we get

$$P(t) = \frac{1}{\lambda + \mu} \left(\begin{bmatrix} \mu & \lambda \\ \mu & \lambda \end{bmatrix} + e^{-(\mu + \lambda)t} \begin{bmatrix} \lambda & -\lambda \\ -\mu & \mu \end{bmatrix} \right).$$

4.1 Birth and Death Process

Definition 20. A birth-and-death (B&D) process is a TH CTMC X on state space $\mathbb{Z}_{\geq 0}$ where $q_{i,j} = 0$ if $j \notin \{i-1, i+1\}$. Let $\lambda_i := q_{i,i+1}$ for $i \geq 0$, $\mu_i := q_{i,i-1}$ for $i \geq 1$, $\mu_0 := 0$.

X(t) is called the population at time t, λ_i is called the birth rate at population i, and μ_i is called the death rate at population i.

Lemma 37. Let X be a B&D process where X(0) = n. Let T_n be the time to reach state n + 1, i.e., $T_n := \min_{t \ge 0} (X(t) = n + 1)$. Then

$$E(T_n) = \frac{1}{\lambda_n} + \frac{\mu_n}{\lambda_n} E(T_{n-1}) = \frac{1}{\lambda_n} \sum_{i=0}^n \prod_{j=1}^i \frac{\mu_{n-j+1}}{\lambda_{n-j}}.$$
$$Var(T_n) = \frac{1}{\lambda_n (\lambda_n + \mu_n)^2} + \frac{\mu_n}{\lambda_n} Var(T_{i-1}) + \frac{\mu_n}{\lambda_n + \mu_n} (E(T_{n-1}) + E(T_n))^2$$

Proof sketch. Let $I_i = \mathbf{1}$ (next transition goes to state i + 1). Let X_i be the transition time out of state *i*. Then $I_i \sim \text{Bernouilli}(\lambda_i/(\mu_i + \lambda_i)), X_i \sim \text{Expo}(\lambda_i + \mu_i)$, and

$$E(T_i \mid I_i) = E(X_i) + (1 - I_i)(E(T_{i-1}) + E(T_i))),$$

$$Var(T_i \mid I_i) = Var(X_i) + (1 - I_i)(Var(T_{i-1}) + Var(T_i)).$$

CKBE for B&D:

$$\frac{dP_{i,j}(t)}{dt} = \mu_i P_{i-1,j}(t) + \lambda_i P_{i+1,j}(t) - (\lambda_i + \mu_i) P_{i,j}(t)$$

CKFE for B&D:

$$\frac{dP_{i,j}(t)}{dt} = \mu_{j+1}P_{i,j+1}(t) + \lambda_{j-1}P_{i,j-1}(t) - (\lambda_j + \mu_j)P_{i,j}(t).$$

Theorem 38 (Limiting Probabilities). Let X be an irreducible B&D process on state space $D \subseteq \mathbb{Z}_{\geq 0}$ where $0 \in D$. For $n \in D$, let $\alpha_n := \prod_{i=1}^n \frac{\lambda_{i-1}}{\mu_i}$. If $\sum_{i \in D} \alpha_i$ is finite, then $P_i = \alpha_i P_0$, and $P_0 = 1 / \sum_{i \in D} \alpha_i$.

Proof sketch. Use Lemma 34 and add adjacent equations.

5 Renewal Theory

Definition 21. Let $[X_1, X_2, ...]$ be a sequence of IID non-negative randvars, called interarrival times, such that $\Pr(X_1 = 0) < 1$ and $\Pr(X_1 = \infty) = 0$. Let $S_n := \sum_{i=1}^n X_i$ (called arrival times). Let $N(t) := \max_n (S_n \leq t)$. Then N is called a renewal process (note that it is a counting process).

We let F and f denote the CDF and PDF/PMF of X_1 , respectively. We let $F^{(n)}$ and $f^{(n)}$ denote the CDF and PDF/PMF of S_n , respectively.

Let R_i be the reward obtained at time X_i for all $i \ge 1$, where all R_i are independent. Let $R(t) := \sum_{i=1}^{N(t)} R_i$. Then R is called a renewal reward process.

Lemma 39. For all $t \ge 0$, $\Pr(N(t) = \infty) = 0$. $\Pr(\lim_{t\to\infty} N(t) = \infty) = 1$.

Proof. Let $\mu := \mathcal{E}(X_1)$. $\mu > 0$ since $\Pr(X_n = 0) < 1$.

$$\Pr\left(\lim_{t \to \infty} \frac{S_n}{n} = \mu\right) = 1. \qquad (\text{strong law of large numbers})$$
$$N(t) = \infty \iff (\forall n, S_n \le t) \implies \lim_{t \to \infty} \frac{S_n}{n} = 0.$$
$$\Pr(N(\infty) = \infty) = 1 \text{ since } \Pr(X_1 = \infty) = 0.$$

Definition 22. For a renewal process N, let $m_N(t) := E(N(t))$. Then m_N is called the mean-value function of N. (If N is clear from context, we will write m instead of m_N .)

Lemma 40.
$$m(t) = \sum_{n=1}^{\infty} \Pr(S_n \le t) = \sum_{n=1}^{\infty} F^{(n)}(t)$$

Theorem 41. *m* uniquely characterizes *F*.

Lemma 42. m(t) is finite for all t.

Theorem 43 (Renewal equation). When interarrival times are continuous randvars,

$$m(t) = F(t) + \int_0^t m(t-x)f(x)dx.$$

Proof sketch. Let $N'(t) := \max_n \left(\sum_{i=2}^{n+1} X_i \leq t \right)$. Then N and N' are identically distributed and

$$N(t) = \begin{cases} 1 + N'(t - X_1) & \text{if } X_1 \le t \\ 0 & \text{if } X_1 > t \end{cases}.$$

Finally, $m(t) = E(E(N(t) \mid X_1)).$

Corollary 43.1. Let N be a renewal process where interarrival times are distributed Uniform (0, 1). Then for $0 \le t \le 1$, $m(t) = e^t - 1$.

Theorem 44 (Limit theorems). For a renewal process N with $\mu := E(X_1)$,

$$\Pr\left(\lim_{t \to \infty} \frac{N(t)}{t} = \frac{1}{\mu}\right) = 1. \qquad \qquad \lim_{t \to \infty} \frac{m(t)}{t} = \frac{1}{\mu}.$$

Theorem 45 (Limit theorems for rewards). For a renewal process N with rewards $\{R_i : i \in \mathbb{Z}_{\geq 1}\}$, let $\alpha := E(R_1)$ and $\mu := E(X_1)$. Then

$$\Pr\left(\lim_{t \to \infty} \frac{R(t)}{t} = \frac{\alpha}{\mu}\right) = 1. \qquad \qquad \lim_{t \to \infty} \frac{\mathrm{E}(R(t))}{t} = \frac{\alpha}{\mu}.$$

Theorem 46 (Central limit theorem for renewals). For a renewal process N with $\mu := E(X_1)$ and $\sigma^2 := Var(X_1)$, the random variable

$$\lim_{t \to \infty} \frac{N(t) - t/\mu}{\sqrt{t\sigma^2/\mu^3}}$$

tends to the standard normal distribution.

Lemma 47 (Stopping time). Let $X = [X_1, X_2, ...]$ be the sequence of interarrival times for renewal process N. Then N(t) + 1 is a stopping time for X.

 $Proof \ sketch. \ N(t)+1 \leq n \iff S_n > t.$

Definition 23. For a renewal process N with arrival times S_1, S_2, \ldots :

- Let $Y(t) := S_{N(t)+1} t$. Y(t) is called the excess at time t.
- Let $L(t) := t S_{N(t)}$. L(t) is called the remaining life at time t.

Lemma 48. Let N be a renewal process with interarrival times $X = [X_1, X_2, \ldots]$. Then $E(S_{N(t)+1}) = t + E(Y(t)) = E(X_1)(m(t)+1)$.

Proof. N(t) + 1 is a stopping time for X.